Perfect Chiral Quantum Routing

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Quantum Walk formalism

The Chiral Hamiltonian:

Graph G : G(Y.E)

• V set of vertices.

• E set of Edges.

The perfect Router:

 $\begin{aligned} & \mathbb{C} \ |\psi_0
angle$; input state $\begin{aligned} & \mathbb{C} \ |\psi_{f,l}
angle$; output state(s)

 $\Re \langle \psi_0 | \psi_{f,l} \rangle = 0$

The lily Graph

Output layer So The Structure: - Routing layer - Chiral layer - Input layer

S Hamiltonian Representation:

$$\circ H(n,d,\vec{\phi}) = A^{in} + A^C(d,\vec{\phi}) + A^R(n,d,\vec{\phi}) + A^{out}(n)$$

 $-A^{in} = |1\rangle\langle 2| + |2\rangle\langle 1|$

$$-A^{C}(d,\vec{\phi}) = \sum_{k \in C} e^{-i\phi_{k}} |2\rangle \langle k| + e^{i\phi_{k}} |k\rangle \langle 2|$$

$$-A^{R}(n, d, \vec{\phi}) = \sum_{k \in C} \sum_{l \in R, l \neq r} |k\rangle \langle l| + |l\rangle \langle k| + \sum_{k \in C} e^{-i\phi_{k}} |r\rangle \langle k| + e^{i\phi_{k}} |k\rangle \langle r|$$

$$-A^{out}(n) = \sum_{a \in R} \sum_{b \in O} |a\rangle \langle b| + |b\rangle \langle a|$$



The Effective Topology



The root of unity $\sqrt[d]{1}$:

$$\sum_{k=1}^{d} e^{-i\phi_k} = \sum_{k=1}^{d} e^{-i\frac{2\pi k}{d}} = 0$$

- © Constructive interference in $|r\rangle$
- \odot Distructive interference in $|l\rangle$



© The dynamics of the walker involves only the desired output(s)

The Dimensionality Reduction Method

The Krylov Reduction Method:

 $\mathfrak{G}\left\langle\zeta\right|e^{-iHt}\left|\psi_{0}\right\rangle=\sum_{k=0}^{\infty}\frac{(-it)^{k}}{k!}\left\langle\zeta\right|H^{k}\left|\psi_{0}\right\rangle=\left\langle\zeta\right|e^{-iH_{red}t}\left|\psi_{0-red}\right\rangle$

 $\mathfrak{G}\mathcal{J}(H,|\zeta\rangle) = span(\{H^k|\zeta\rangle|k \in \mathbb{N}_0\})$



🛞 Reduced Ba*s*is:

$$\circ \quad \left| e_1 \right\rangle = \left| f \right\rangle$$

$$\circ |e_2\rangle = |r\rangle$$

$$\circ ||e_3\rangle = \frac{1}{\sqrt{d}} \sum_C e^{i\phi_k} |k\rangle$$

$$\circ ||e_4\rangle = |2\rangle$$

 $\circ |e_5\rangle = |1\rangle$





S Basis of identically evolving vertices

$$\mathfrak{G} H_{red} = \langle e_j | H | e_k \rangle$$

The Purely Topological Routing Process

$$H_{red}(d) = \begin{pmatrix} 0 & 1 & 0 & 0 & 0 \\ 1 & 0 & \sqrt{d} & 0 & 0 \\ 0 & \sqrt{d} & 0 & \sqrt{d} & 0 \\ 0 & 0 & \sqrt{d} & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 \end{pmatrix}$$

 $\$ Input State: $|\psi_0
angle = lpha |1
angle + e^{i\gamma}\sqrt{1-lpha^2} |2
angle$

 $\circ \ lpha = 0$; lpha = 1 routing of classical information

figs Output State: $\left|\psi_{f}
ight
angle=lpha|f
angle+e^{i\gamma}\sqrt{1-lpha^{2}}|r
angle$

$$P_{1f}(t) = \frac{(2d - (2d + 1)\cos(t) + \cos(t\sqrt{1 + 2d}))^2}{4(1 + 2d)^2}$$
$$P_{2r}(t) = \frac{(\cos(t) - \cos(t\sqrt{1 + 2d}))^2}{4}$$



Fig: Optimal time t^* (in unit of π) maximizing the routing probability $P_{1f}(t^*)$ as a function of the chiral layer dimension d.



Fig: Logplot of the maximized routing probability $P_{1f}(t^*)$ as a function of the chiral layer dimension d (red). The difference $|P_{1f}(t^*) - P_{2r}(t^*)|$ (blue).



<u>Nearly</u> optimal Chiral Quantum routing

The Weighted Routing Process

$$\circ H(n,d,\vec{\phi}) = A^{in} + \beta A^{C}(d,\vec{\phi}) + \beta A^{R}(n,d,\vec{\phi}) + A^{out}(n)$$

$$H_{red,\beta}(\beta,d) = \begin{pmatrix} 0 & 1 & 0 & 0 & 0 \\ 1 & 0 & \beta\sqrt{d} & 0 & 0 \\ 0 & \beta\sqrt{d} & 0 & \beta\sqrt{d} & 0 \\ 0 & 0 & \beta\sqrt{d} & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 \end{pmatrix}$$

$$\begin{split} & \text{(S) Input State: } |\psi_0\rangle = \alpha |1\rangle + e^{i\gamma}\sqrt{1-\alpha^2}|2\rangle \\ & \text{(S) Output State: } |\psi_f\rangle = \alpha |f\rangle + e^{i\gamma}\sqrt{1-\alpha^2}|r\rangle \end{aligned}$$

 $\circ \ lpha = 0$; lpha = 1 routing of classical information

$$P_{1f}(t) = \frac{(2d\beta^2 - (2d + \beta^2)\cos(t) + \cos(t\sqrt{\beta^2 + 2d}))^2}{4(\beta^2 + 2d)^2}$$

$$P_{2r}(t) = \frac{(\cos(t) - \cos(t\sqrt{\beta^2 + 2d}))^2}{4}$$

Robustness & Time Periodicity

Taylor Expansion of the Probabilities:

$$P_{1f}(t) = 1 - (t - t_{\beta}^{*})^{2} + O(t - t_{\beta}^{*})^{3}$$
$$P_{2r}(t) = 1 - \frac{5}{2}(t - t_{\beta}^{*})^{2} + O(t - t_{\beta}^{*})^{3}$$



Extension to Qudit Routing and Physical Implementations





$$\mathfrak{S} e^{-i\widetilde{H}(2q+1)\pi} = \sum_{s=0}^{2\widetilde{D}} |e_{1+s}\rangle \langle e_{2\widetilde{D}+1-s}|, q \in \mathbb{N}$$

$$= \begin{pmatrix} 0 & \cdots & 0 & 1 \\ \vdots & \ddots & \ddots & 1 & 0 \\ \vdots & \ddots & \ddots & 0 \\ 0 & 1 & \cdots & 0 \end{pmatrix}$$

Sinput State: $|\psi_0\rangle = \sum_{l=1}^{\widetilde{D}} \alpha_l |l\rangle$

SOutput State:
$$|\psi_f\rangle = \alpha_{\widetilde{D}}|r\rangle + \sum_{l=0}^{\widetilde{D}-1} \alpha_l |f_{2\widetilde{D}+1-l}\rangle$$

Perfect Periodic Chiral Quantum routing

Summary & Conclusions

- $\ensuremath{\mathbb{C}}$ Nearly perfect quantum routing only with topology and chirality
- © Perfect Chiral quantum routing with the weighted configuration



© Both for quantum and classical Information (bit, qubit and qudit)

© Universality of the time required promote our result to the most general routing procedure